

2009c

H

From verses to stacks in 2 hours

$$\begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} \otimes \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$$

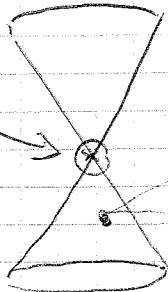
$$X = C^2$$

$$C = \mathbb{Z}/2\mathbb{Z}$$

quiet stack

$$[X/C]$$

$C[X] \otimes C[X]$
quiet variety



$$X/C$$

$$x^2, y^2, xy$$

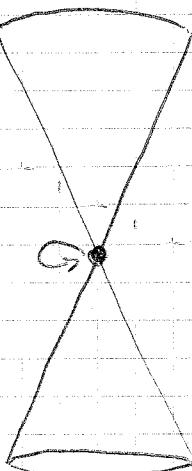
$$z^2$$

$$W(\mathbb{Z}/2)$$

$$X/C$$

nonplanar
reason: higher stack non planar

extra info about standard
nugget.

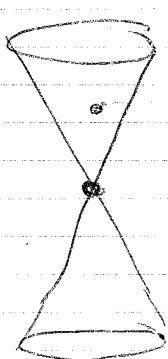


known smooth as stack

noncommutative crossed
 $C[X] \# C$ smash product
 $(f \# g)(f' \# g') = f f' \# g g'$
cate = $C[X \# C]$



non non
 $C[X] \# C$



cate

pepsi app
of one square

$C[X] \# C$ is "smooth"

oldie & do + max

$$C^3$$

Calabi-Yau orbifolds

spaced:

kleinian nifulari

pepsi app
of one square

cate

CATEGORY THEORY 101

A category consists of

- 1. objects
- 2. morphism btw objects
allowing composition and
identity morph.

Example ① (f, g algs) has

objects: all f, g form \mathbb{C} -algebra

$$C[X] = \frac{(tx_1 \rightarrow x_n)}{(f_1 \rightarrow f_n)}$$

morphism: all \mathbb{C} -algebra morphism

$$C[X] \xrightarrow{\cong} C[Y]$$

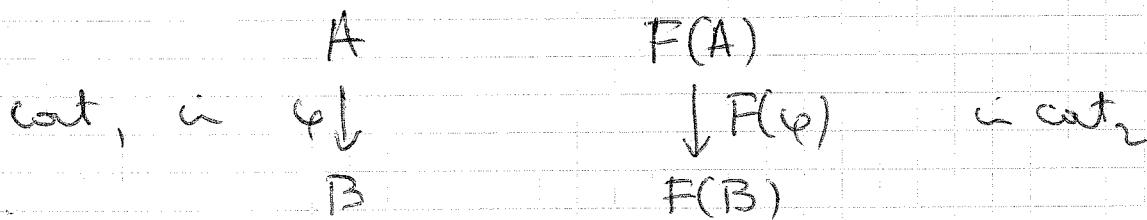
② (sets) have objects sets
morph maps.

③ (groups) object group
morph groupmorphe.

A FUNCTOR is a 'map' between categories

$$F: (\text{cat}_1) \rightarrow (\text{cat}_2)$$

sending objects to objects and morphism
to morphism in a compatible way.



Examp: $\odot \oplus_m : (\text{fg algs}) \rightarrow (\text{groups})$

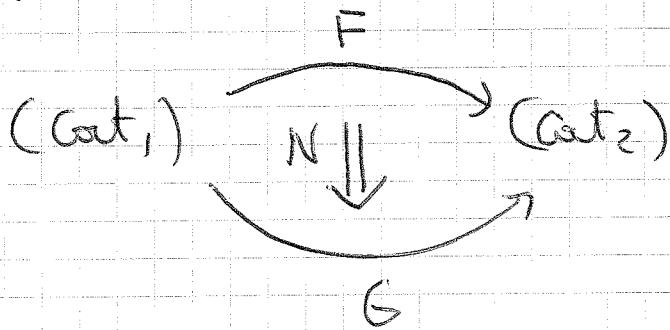
10

$\mathbb{C}[x] \mapsto (\mathbb{C}[x])^*$ with \odot

④ GL_n : $(\text{fg algs}) \rightarrow (\text{groups})$

$\mathbb{C}[x] \mapsto GL_n(\mathbb{C}[x])$

A NATURAL TRANSFORMATION is a map between two functors btw some categories



such a map $\in (\text{Cat}_2)$ so every object $\in (\text{Cat}_1)$

$$F(A) \xrightarrow{\quad} G(A)$$

$$N(A)$$

such that for all maps $A \xrightarrow{\varphi} B$ in (Cat_1) the
comute diagram in (Cat_2)

$$\begin{array}{ccc} F(A) & \xrightarrow{N(A)} & G(A) \\ F(\varphi) \downarrow & & \downarrow G(\varphi) \\ F(B) & \xrightarrow{N(B)} & G(B) \end{array}$$

Examp: obj: $GL_n \Rightarrow \oplus_m$ is natural trans.

$$\left\{ \begin{array}{l} p_1(x_1, \dots, x_n) = 0 \\ \vdots \\ p_k(x_1, \dots, x_n) = 0 \end{array} \right.$$

$$I = (p_1, \dots, p_k) \subset \mathbb{C}[x_1, \dots, x_n]$$

④

$$X = V(I)$$

$$p_1(x_1, \dots, x_n) = 0$$

$$\mathbb{C}[x] = \frac{\mathbb{C}[x_1, \dots, x_n]}{(I)}$$

$V(I) \hookrightarrow \{ \text{max ideals of } \mathbb{C}[x] \}$

"alg maps $\mathbb{C}[x] \rightarrow \mathbb{C}$ "

(affine)

$$\stackrel{\cong}{\longrightarrow}$$

(fg algebras)

objects: zero set

$$V(p_1, \dots, p_k) \subset \mathbb{A}^n$$

"
X

objects: f.g. \mathbb{C} -algebra

$$(p_1, \dots, p_k) \rightarrow (\mathbb{C}[x_1, \dots, x_n]) \rightarrow \mathbb{R}$$

$$\mathbb{C}[x] = R \cong \frac{\mathbb{C}[x_1, \dots, x_n]}{(p_1, \dots, p_k)}$$

morphism: \mathbb{C} -algebra morph.

$$X \leftarrow Y$$



$$\mathbb{C}[x] \rightarrow \mathbb{C}[y]$$

c

problem: Hilbert Nullstellensatz

different algebras same zero set

$$\text{ex: } \underline{\mathbb{C}[x_1, \dots, x_n]}$$

$$(p(x_1, \dots, x_n))$$

$$\underline{\mathbb{C}[x_1, \dots, x_n]}$$

$$(p(x_1, \dots, x_n)^e)$$

have

and
less

then is only restriction.

If we want $(\text{affine}) \cong (\text{fg-algs})^\circ$ have to add
extra data to zero set X

3 approach

(Serre) : On X put "structure sheaf" \mathcal{O}_X
having global sections $\mathcal{C}(X) = \Gamma(X, \mathcal{O}_X)$

(Borel-Weil-Bott) : "functor of points"

$h_X : (\text{affine}) \rightarrow (\text{sets})$ contravariant for
 $y \mapsto \text{Hom}(Y, X)$
 $= \text{Hom}_{\text{alg}}(\mathcal{C}(X), \mathcal{C}(Y))$

$h_X(C) = \text{Hom}_{\text{alg}}(\mathcal{C}(X), C) = \text{pts of } X$

$h_X(\frac{C(t)}{(t^2)}) = \text{Hom}_{\text{alg}}(\mathcal{C}(X), \frac{C(t)}{(t^2)}) = \text{pts + tangent vector}$
 $= \text{pts of } TX$

⋮

(affine) objects: functors h_X
 correspond to f, g alg $\mathcal{C}(X)$

How to relate with } Morphism should be
 morphism in } $h_X \rightarrow h_Y$ natural
 $(fg\text{-alg})$? } transformation.

YONEDA

$\text{Hom}(h_X, h_Y) = h_Y(X) = \text{Hom}_{\text{alg}}(\mathcal{C}(X), \mathcal{C}(Y))$

$\uparrow \dashv$

mod $\Leftrightarrow \mathbb{C}[X] \xrightarrow{\epsilon} \mathbb{C}[Y]$

(3)

$$\begin{array}{ccc} & \swarrow & \downarrow f \\ \mathbb{C}[Z] & & \end{array}$$

$\nexists \alpha(z) \in \text{Ob}(fg_{\text{alg}})$

$$\text{Hom}_{\text{alg}}(\mathbb{C}[Z], \mathbb{C}[X]) \rightarrow \text{Hom}_{\text{alg}}(\mathbb{C}[Z], \mathbb{C}[Y])$$

||

||

$$h_X(z) \rightarrow h_Y(z)$$

so φ determines function $h_X \rightarrow h_Y$
nat. transfo

\Rightarrow given natural transfo

$$h_X \xrightarrow{N} h_Y$$

$$\forall z : h_X(z) \xrightarrow{N(z)} h_Y(z)$$

in particular

$$id_X \in h_X(X) \rightarrow N(X) | id_X \in h_Y(X) = \text{Hom}(X, Y)$$

□

(Kontsevich) : dual coalgebra

take as geometric object

$$\mathbb{C}[X]^{\circ} = \{ \lambda : \mathbb{C}[X] \xrightarrow{\text{lin}} \mathbb{C} \mid K(\lambda) \text{ contains coprime } \text{(de)} \}$$

is coalgebra

coradical recovers pts of X and $\mathbb{C}[X]^{\circ} = \prod_{x \in X} \hat{\mathbb{C}}_{x,x}$

Frobenius in artinian geometry

what is special about function

$$h_x : (\text{affine}) \rightarrow (\text{sets})$$

They are SHEAVES for the $\overline{\text{ETALE}}$ TOPOLOGY

Etale extension $y \xrightarrow{\text{et}} x$ iff.

$$\mathcal{O}[x] \rightarrow \mathcal{O}[y] = \mathcal{O}[x] \frac{[z_1, \dots, z_n]}{(f_1, \dots, f_k)}$$

a direct sum of these

Etale morphism
requires
for failure of
given f_i or the
 n alg geo

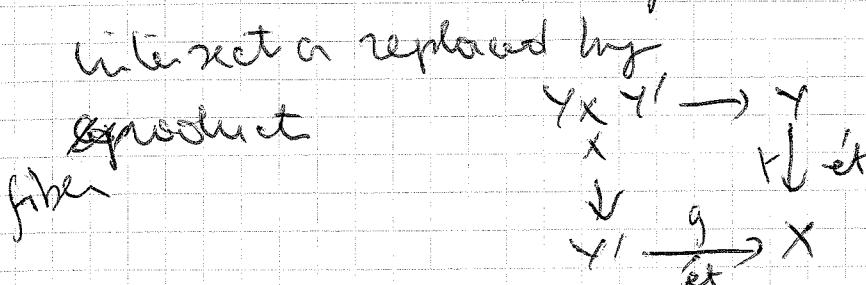
$$\det \begin{bmatrix} \frac{\partial f_1}{\partial z_1} & \cdots & \frac{\partial f_1}{\partial z_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_k}{\partial z_1} & \cdots & \frac{\partial f_k}{\partial z_n} \end{bmatrix} \in \mathcal{O}[y]^*$$

Basic example

$$\mathcal{O}[x] \rightarrow \mathcal{O}[x] \frac{[z]}{S(z^n - a)} \quad a \in \mathcal{O}[x]^*$$

should think of $y \xrightarrow{\text{et}} x$ as substitute for
Zariski open $U \subset X$

Etale for Frobenius
topology
intersect a replaced by
use intersection.



$$Y \times Y' = \{ (y, y') : f(y) = g(y') \text{ in } X \}$$

$$\mathbb{C}[Y \times Y'] = \mathbb{C}[X] \otimes \mathbb{C}[Y']$$

$$= \mathbb{C}[X] \frac{[z_1 \rightarrow z_2]}{(f_1, \rightarrow f_2)} \otimes \mathbb{C}[X] \frac{[v_1 \rightarrow v_2]}{(g_1, \rightarrow g_2)}$$

$$= \mathbb{C}[X] \frac{[z_1 \rightarrow z_2, v_1 \rightarrow v_2]}{(f_1, \rightarrow f_2, g_1, \rightarrow g_2)} \quad \text{Jac} = \begin{bmatrix} * & 0 \\ 0 & * \end{bmatrix}$$

so this is étale.

$$\begin{array}{ccc} Y \times Y' & \xrightarrow{\text{ét}} & Y \\ \times & \downarrow \text{ét} & \downarrow \text{ét} \\ Y' & \xrightarrow{\text{ét}} & X \end{array}$$

Blade cover is étale map surjective on points.

Example $X = A^n$ $\mathbb{C}[X] = \mathbb{C}[x_1, \rightarrow x_n]$ has only units \mathbb{C}^*

so has no global étale maps

$$\mathbb{C}[x_1, \rightarrow x_n][z] = \mathbb{C}[x_1, \rightarrow x_n]$$

$$\begin{array}{c} A^n \\ A^n \\ A^n \\ \vdots \\ A^n \end{array}$$

But can look at étale maps over affine open subsets

$$\mathbb{C}(x_{-i}) \quad \mathbb{C}(y_i) = \mathbb{C}[x, x^{-1}, y] \frac{[z]}{(z^2 - x)} \quad \text{is étale ext of } \mathbb{C}(x, x^i, y)$$

$$\begin{array}{ccc} \mathbb{A}(x) & \xrightarrow{\quad} & \mathbb{A}(x, (x-1)^{-1}, y) \\ \downarrow & & \downarrow \\ \mathbb{A}(x) & & \mathbb{A}(x, (x-1)^{-1}, y) \end{array}$$

$$Y = Y_1 \cup Y_2$$

$$\mathbb{C}[Y] = \mathbb{C}[Y_1] \oplus \mathbb{C}[Y_2]$$

is étale et of (\mathbb{A}^1, y)

so there are plenty of étale covers of affine schemes
and these covers can be refined by taking ^{like} ~~products~~

SHEAF on ETALÉ TOPOLOGY

Thm (Grothendieck descent) $Y \xrightarrow{\text{ét}} X$ étale cover

Then for all $Z \in \text{Ob}(\text{affine})$

$$\text{Hom}(X, Z) \longrightarrow \text{Hom}(Y, Z) \xrightarrow{\text{ii}} \text{Hom}(Y \times^Y_X Z, Z)$$

$$\begin{array}{ccc} Y \times^Y_X Z & \longrightarrow & Y \\ \downarrow & \nearrow & \downarrow \\ Z & \longrightarrow & X \end{array}$$

We say that \mathcal{H}_Z has a sheaf on (affine)

So far: only considered affine varieties

want to "glue" them together to get more

interesting geometrical objects (like projective

varieties, moduli spaces etc.)

7

Def A SPACE is a contravariant functor

$$f : (\text{affine}) \rightarrow (\text{sets})$$

enjoying the sheaf property for the étale topology
 i.e. if $Y \xrightarrow{\text{ét}} X$ étale covers in (affine)
 have

$$f(X) \rightarrow f(Y) \rightrightarrows f\left(\underset{X}{Y \times Y}\right)$$

This is still too loose a condition. We want to restrict to spaces admitting an ATLAS like in diff. geometry but now with charts \hookrightarrow affine schemes.

Space f has a set of points $|f| = f(C)$ and want to cover them with (points) of affine sets.

$$U = U_1 \sqcup U_2 \sqcup \dots \sqcup U_n \longrightarrow |f|$$

$$C[U] = (U_1) \oplus \dots \oplus (U_n)$$

but need to say when points in different charts are mapped to the same point in the space f

$$\begin{array}{ccc} R = U \times U & \xrightarrow{m} & R \text{ closed subspace of} \\ \downarrow f_1 & & \downarrow \\ U & \xrightarrow{p_2} & R \text{ is equivalence relation.} \end{array}$$

defn an ALGEBRAIC SPACE \mathcal{F} is a SPACE \mathcal{F} having an étale ATLAS

That is, \exists affine scheme U and closed subspace

$$R \subset U \times U \quad (\text{i.e. } \mathcal{O}(R) \subset \mathcal{O}(U) \otimes \mathcal{O}(U))$$

s.t. $\begin{cases} \text{① } R \text{ is equivalence relation} \\ \text{② two projections } R \xrightarrow{\text{pr}_i} U \text{ are étale monomorphisms} \\ \text{③ } |f| = \frac{|U|}{|R|} \end{cases}$

is étale atlas.

defn a SCHEME is an ALGEBRAIC SPACE

such that the restriction of R to each connected component of U is the trivial diagonal equivalence relation.

CATEGORY THEORY 201

[9]

A 2-CATEGORY is a category such that
for all $A, B \in \text{Ob}(B)$ $\text{Hom}(A, B)$ is a category

That is, B consists of

- Objects
- Morphisms (which are objects in a cat)
- 2-morphisms (~~morph~~^{morph} between morphs satisfying "natural" condition as elements in the cat)

Example • (cat) { objects: categories
 { morph: functors
 2-morph: natural transforms

- (groupoids)
 - objects: groupoids (i.e. cat. in which all morphs are isompr.)
 - morphism: functors
 - 2-morph: natural transfor.

• Every category B becomes 2-category in obvious way: make $\text{Hom}(A, B)$ a category with objects the elements of $\text{Hom}(A, B)$ every object has identity map (and no others)

• (affine) as 2-category

• (affine/ X) as 2-category

↑ all Y with "stochastic maps" $Y \rightarrow X$

Recall: Space is contravariant functor

$\mathcal{F} : (\text{affine}) \rightarrow (\text{sets})$

which is sheaf in étale top on (affine)

Dfn

a STACK is a contravariant 2-functor

$\mathcal{F} : (\text{affine}) \rightarrow (\text{groupoids})$

as 2-set

what is a sheaf in étale top on (affine).

So: what is a 2-functor?

sends
objects $\begin{cases} \text{objects} \rightarrow \text{objects} \\ \text{Morph} \rightarrow \text{morphism} \\ 2\text{-morph} \rightarrow 2\text{-morphism} \end{cases}$

So: a stack associates to an affine scheme X

a groupoid $\mathcal{F}(X)$

To a morphism of affine schemes $X \rightarrow Y$ (an alg. map $\mathcal{O}[Y] \rightarrow \mathcal{O}[X]$) a functor

$$\mathcal{F}(X) \leftarrow \mathcal{F}(Y)$$

and to each 2-morphism in (affine)

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ & \Downarrow id & \\ & \xleftarrow{g} & \end{array}$$

a natural transfo

$$\begin{array}{ccc} \mathcal{F}(X) & \xrightarrow{\mathcal{F}(f)} & \mathcal{F}(Y) \\ \Downarrow \mathcal{F}(id) & & \Downarrow \mathcal{F}(g) \\ \mathcal{F}(f) & & \mathcal{F}(g) \end{array}$$

but note: this does not have to be the identity natural — only natural is 0.

Example X affine G finite gp acting on X (1)
quotient stack

$[X/G] : (\text{affine}) \rightarrow (\text{groupoids})$ containing
variant

$$Y \mapsto [X/G](Y)$$

$[X/G](Y)$: objects are triples

$$\begin{array}{ccc} E & \xrightarrow{f} & X \\ \pi \downarrow & & \\ Y & & \end{array}$$

π defines principal
 G -bundle \hookrightarrow étale top

f is G -equivariant.

morphisms in $[X/G](Y)$ are

$$\begin{array}{ccc} & X & \\ f \nearrow & \swarrow f' & \\ E & \dashrightarrow & E' \\ \pi \searrow & \swarrow \varphi & \pi' \\ Y & & \end{array}$$

where φ must be
 G -equiv. isomorph.

Morphisms actisted or by $[X/G]$

gives functor $[X/G](+)$

$$Y' \xrightarrow{+} Y \quad [X/G](Y) \rightarrow [X/G](Y')$$

$$\begin{array}{ccccc} Y \times \mathbb{R} & \xrightarrow{\exists!} & E' & \dashrightarrow & E \\ \downarrow & & \downarrow & & \downarrow \\ Y' & \longrightarrow & Y & & \end{array}$$

2-morphism becomes

Demystification

$$X = \mathbb{A}^2$$

$$\mathcal{O}(X) = \mathcal{O}(x, y)$$

(12)

$$G = \mathbb{Z}/2\mathbb{Z}$$

action: point-reflection in origin



what is $[X/G](Y)$?

$\mathbb{E} \xrightarrow{\text{forget}} X$ (i) G -torsors are precisely
 $\begin{array}{c} \pi \\ \downarrow \\ Y \end{array}$

$$\mathcal{O}[E] = \mathcal{O}[Y] \frac{[t]}{(t^2 - a)} \quad a \in \mathcal{O}[Y]^*$$

so are in particular étale

(ii) $\mathcal{O}[X] \rightarrow \mathcal{O}[E] = \mathcal{O}[Y] + \mathcal{O}[Y]\sqrt{a}$

$$\begin{cases} x \mapsto \lambda\sqrt{a} \\ y \mapsto \mu\sqrt{a} \end{cases}$$

What are G -morphisms

$\begin{array}{ccc} \mathbb{E} & \xrightarrow{\quad X \quad} & E' \\ \downarrow & \nearrow \varphi & \downarrow \\ Y & \xrightarrow{\quad b \quad} & Y' \end{array}$

 $\varphi: \mathcal{O}[E'] \rightarrow \mathcal{O}[E]$

$$(\mathcal{O}[Y] \oplus \mathcal{O}[Y]\sqrt{a}) \xrightarrow{\quad \text{?} \quad} (\mathcal{O}[Y] + \mathcal{O}[Y]\sqrt{a})$$

$\varphi \in \mathcal{O}[Y]$ -morphism so

$$\sqrt{a} \mapsto b\sqrt{a}$$

so that

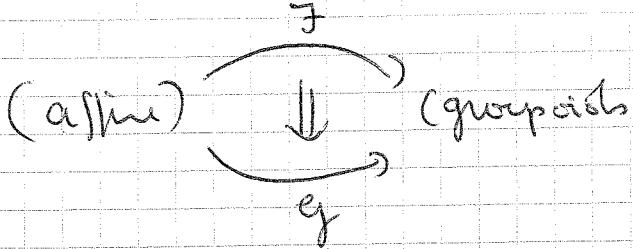
$$\varphi: a' = b^2 a$$

$$\varphi: a = b^2 a'$$

so $b \in \mathcal{O}[Y]^*$

so torsor

Morphisms between stacks are natural transformations



YONEDA for stacks

X affine \mathcal{F} stack
as stack

$$\text{Hom}(X, \mathcal{F}) \cong \mathcal{F}(X)$$

Proof: " $X \xrightarrow{\cong} \mathcal{F}$ " $\Leftrightarrow N\left(\frac{X}{\mathcal{F}}\right) = 0$

Importance: Can view elements of $\mathcal{F}(X)$ as maps $X \rightarrow \mathcal{F}$

Example: There is a special element in $[X/G](X)$...

$$\begin{array}{ccc} X \times G & \xrightarrow{\text{action}} & X \\ \pi_1 \downarrow & & \downarrow u \\ X & & [X/G] \end{array}$$

by Yoneda corresponds to map

MAIN THEOREM ON QUOTIENTS STACKS

$X \xrightarrow{u} [X/G]$ is representable surjective étale map giving
etale for the stack which the becomes an
"algebraic" stack. Moreover X is a universal
family for the quotient problem.

X affine, Z stack

$\Omega^0 X \rightarrow Z$ is [representable] iff $\forall Y$ affine $\forall Y \rightarrow Z$

$$\begin{array}{ccc} U & \cong & Y \times X \rightarrow X \\ & & \downarrow \quad \downarrow \\ & & Y \rightarrow Z \end{array}$$

from like product

$Y \times X$ which is stack

(affine) \longrightarrow (affine + étale)

$$\begin{array}{ccc} Z & \xrightarrow{\quad} & X \\ \downarrow & & \downarrow \\ Y & \xrightarrow{\quad} & Z \end{array}$$

obj₁

obj₂

$Z \xrightarrow{\quad} \{Z, \alpha\} : \alpha \in \alpha$

with structure
maps to $X \otimes Y$

$Z(Z)$ between

objects obtained by

(affine)

This fibre-product stack is representable by a scheme U

e) when map $X \rightarrow Z$ is representable the properties of the maps hold if and only then property holds for all maps $U \rightarrow Y$ coming from representability

e.g. surjective means all $U \rightarrow Y$ surject

étale

"

étale etc.

PROOF of main thm: claim:

E

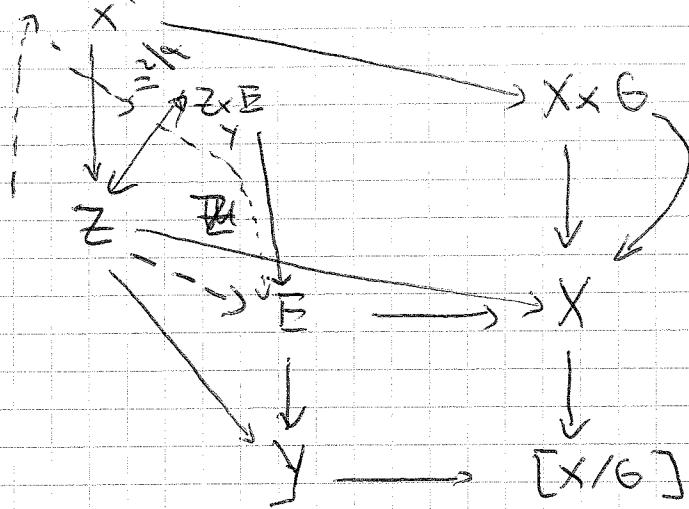
$$\begin{array}{ccc} E & \xrightarrow{\text{epi}} & X \\ \downarrow \text{perf} & & \downarrow \text{det} \\ Y & \xrightarrow{\quad} & [X/G] \end{array}$$

$$Y \times X \cong E$$

$[X/G]$

and no map will be inj + étale
because $E \rightarrow Y$ is.

$$\mathbb{Z} \times (X \times G) \cong \mathbb{Z} \times G$$



□

Example: $X = \mathbb{A}^2$

with $G = \mathbb{C}/\mathbb{C}^\times$ action

~~$\mathbb{C}/\mathbb{C}^\times$ acts on \mathbb{A}^2~~

What are points of $[X/G]$, that is, $[X/G](\mathbb{C})$

There is only one G -fiber over 0 in $\mathbb{C} \times \mathbb{C} = \frac{\mathbb{C} \times \mathbb{C}}{(\mathbb{C}^\times)^2}$
 with action $(\mathbb{C} \times \mathbb{C})$. So element of $[X/G](\mathbb{C})$ is connect algebraic closure

$$\begin{array}{ccc}
 & x \vdash (\alpha, \beta) & \\
 \mathbb{C} \times \mathbb{C} & \xleftarrow{\varphi} & \mathbb{C}[x,y] \\
 \uparrow & & \\
 \mathbb{C} & &
 \end{array}$$

$x \vdash (\alpha, \beta)$
 $-x \vdash (-\alpha, -\beta) = (\beta, \alpha)$
 $\text{so } \begin{cases} x \vdash (\alpha, -\alpha) \\ y \vdash (\beta, -\beta) \end{cases}$ one only possibility

That is $\mathbb{C}[x,y] \xrightarrow{\varphi} \mathbb{C} \times \mathbb{C} \xrightarrow{m} \mathbb{C}$ determine the
 $\downarrow p_m$ two points $(\alpha, \beta), (-\alpha, -\beta) \in \mathbb{A}^2$

so elements of $[X/G]$ are precisely the orbits.

But: not just orbits, also have to compute actions in $[X/G](\mathbb{C})$

$$\begin{array}{c} \mathbb{C}[X,Y] \\ \varphi \swarrow \quad \bullet \quad \searrow \varphi \\ \mathbb{C} \times \mathbb{C} \dashrightarrow \mathbb{C} \times \mathbb{C} \\ \text{C} \end{array}$$

If $\alpha \neq 0$ then β non-trivial automorph in \mathbb{C}^*

If $\alpha = 0$ we have $\left\{ \begin{array}{l} \text{ct} \\ \text{twist} \end{array} \right\}$ as automorph.

So $[X/G](\mathbb{C})$ is

